

IN THE CLAIMS:

1 1. (Currently Amended) A method for executing uniprocessor (UP) coded workloads in a
2 multiprocessor (MP) computer system without having to rewrite the UP-coded work-
3 loads' code, the method comprising the steps:

4 organizing the UP-coded workloads into one or more concurrency groups,
5 wherein UP-coded workloads in the same concurrency group are not permitted to execute
6 concurrently with one another in the MP computer system;

7 scheduling first and second execution vehicles that respectively execute on differ-
8 ent processors in the MP computer system at substantially the same time;

9 acquiring a first concurrency group by the first execution vehicle and a second
10 concurrency group by the second execution vehicle; and

11 executing UP-coded workloads in the first concurrency group through the first
12 execution vehicle at substantially the same time as UP-coded workloads in the second
13 concurrency group are executed through the second execution vehicle.

1 2. (Original) The method according to claim 1, wherein the UP-coded workloads are UP-
2 coded threads, and the first and second execution vehicles are first and second processes.

1 3. (Original) The method according to claim 1, wherein the UP-coded workloads are
2 messages, and the first and second execution vehicles are first and second threads.

1 4. (Original) The method according to claim 1, wherein the step of acquiring the first and
2 second concurrency groups further comprises:

3 dequeuing from a concurrency-group run queue a first concurrency-group data
4 structure associated with the first concurrency group; and

5 dequeueing from the concurrency-group run queue a second concurrency-group
6 data structure associated with the second concurrency group.

1 5. (Original) The method according to claim 4, further comprising:

2 setting a first CG flag in the first concurrency-group data structure to a value indi-
3 cating that the first concurrency group is in a running state; and

4 setting a second CG flag in the second concurrency-group data structure to a
5 value indicating that the second concurrency group is in a running state.

1 6. (Original) The method according to claim 4, further comprising:

2 appending UP-coded workloads enqueued on a first current queue in the first con-
3 currency-group data structure onto a first active queue in the first concurrency-group data
4 structure; and

5 appending UP-coded workloads enqueued on a second current queue in the sec-
6 ond concurrency-group data structure onto a second active queue in the second concur-
7 rency-group data structure.

1 7. (Original) The method according to claim 6, further comprising:

2 dequeueing UP-coded workloads in the first and second concurrency groups from
3 the first and second active queues, respectively; and

4 executing the dequeued UP-coded workloads to completion.

1 8. (Original) The method according to claim 5, further comprising:

2 in response to the first execution vehicle finishing execution of the UP-coded
3 workloads in the first concurrency group, the first execution vehicle performing the steps:

4 A) if at least one UP-coded workload in the first concurrency group is
5 executable:

6 (i) setting the value of the first CG flag to a value indicat-
7 ing that the first concurrency group is in a queued state;

- 8 (ii) re-enqueueing the first concurrency-group data struc-
9 ture onto the concurrency-group run queue;
- 10 B) if there are not any UP-coded workloads in the first concurrency
11 group that are executable, setting the first CG flag to a value indicating that the
12 first concurrency group is in a suspended state;
- 13 C) dequeuing from the concurrency-group run queue a third concur-
14 rency-group data structure associated with a third concurrency group; and
- 15 D) setting a third CG flag in the third concurrency-group data structure to
16 a value indicating that the third concurrency group is in a running state.
- 1 9. (Original) The method according to claim 1, wherein at least one of the UP-coded
2 workloads is organized into the one or more concurrency groups at run-time.
- 1 10. (Original) The method according to claim 1, wherein the MP computer system is a
2 network cache.
- 1 11. (Original) A multiprocessor (MP) computer system configured to execute uniproc-
2 sor (UP) coded threads without having to rewrite the UP-coded threads' code, the MP
3 computer system comprising:
4 a plurality of processors;
5 a memory having a plurality of storage locations addressable by the plurality of
6 processors for storing data and program code, the memory being configured to store a
7 separate concurrency-group data structure for each of a plurality of concurrency groups,
8 each concurrency-group data structure comprising:
9 an active-queue pointer storing a location in the memory of an active
10 queue of UP-coded thread messages associated with UP-coded threads in an ex-
11 ecutable state; and

12 a current-queue pointer storing a location in the memory of a current
13 queue of UP-coded thread messages associated with UP-coded threads waiting to
14 be transferred to the active queue.

1 12. (Original) The MP computer system according to claim 11, wherein each concur-
2 rency-group data structure further comprises a CG flag that stores a value indicating an
3 operational state of a concurrency group associated with the concurrency-group data
4 structure.

1 13. (Original) The MP computer system according to claim 11, wherein each UP-coded
2 thread message stored in the active queue and current queue stores a location in the
3 memory of a top of a call stack associated with a specific UP-coded thread.

1 14. (Original) The MP computer system according to claim 13, wherein the call stack is
2 accessible through a thread control block (TCB) associated with the specific UP-coded
3 thread, the TCB including a CG pointer for storing a memory location of a concurrency-
4 group data structure.

1 15. (Original) The MP computer system according to claim 11, wherein each concur-
2 rency-group data structure further comprises meta-data information associated with a
3 concurrency group.

1 16. (Original) The MP computer system according to claim 11, wherein the MP computer
2 system is a network cache.

1 17. (Currently Amended) An apparatus for executing uniprocessor (UP) coded workloads
2 in a multiprocessor (MP) computer system without having to rewrite the UP-coded work-
3 loads' code, the method comprising the steps:

means for organizing the UP-coded workloads into one or more concurrency groups, wherein UP-coded workloads in the same concurrency group are not permitted to execute concurrently with one another in the MP computer system;

means for scheduling first and second execution vehicles that respectively execute on different processors in the MP computer system at substantially the same time;

means for acquiring a first concurrency group by the first execution vehicle;

means for acquiring a second concurrency group by the second execution vehicle;

and

means for executing UP-coded workloads in the first concurrency group through the first execution vehicle at substantially the same time as UP-coded workloads in the second concurrency group are executed through the second execution vehicle.

18. (Original) The apparatus according to claim 17, wherein the UP-coded workloads are UP-coded threads, and the first and second execution vehicles are first and second processes.

19. (Original) The apparatus according to claim 17, wherein the UP-coded workloads are messages, and the first and second execution vehicles are first and second threads.

20. (Original) The apparatus according to claim 17, further comprising:

means for dequeuing from a concurrency-group run queue a first concurrency-group data structure associated with the first concurrency group; and

means for dequeuing from the concurrency-group run queue a second concurrency-group data structure associated with the second concurrency group.

21. (Original) The apparatus according to claim 20, further comprising:

means for setting a first CG flag in the first concurrency-group data structure to a value indicating that the first concurrency group is in a running state; and

4 means for setting a second CG flag in the second concurrency-group data struc-
5 ture to a value indicating that the second concurrency group is in a running state.

1 22. (Original) The apparatus according to claim 20, further comprising:

2 means for appending UP-coded workloads enqueued on a first current queue in
3 the first concurrency-group data structure onto a first active queue in the first concur-
4 rency-group data structure; and

5 means for appending UP-coded workloads enqueued on a second current queue in
6 the second concurrency-group data structure onto a second active queue in the second
7 concurrency-group data structure.

1 23. (Original) The apparatus according to claim 22, further comprising:

2 means for dequeuing UP-coded workloads in the first and second concurrency
3 groups from the first and second active queues, respectively; and

4 means for executing the dequeued UP-coded workloads to completion.

1 24. (Original) The apparatus according to claim 21, further comprising:

2 means for setting the value of the first CG flag to a value indicating that the first
3 concurrency group is in a queued state or in a suspended state; and

4 means for re-enqueueing the first concurrency-group data structure onto the con-
5 currency-group run queue.

1 25. (Currently Amended) A computer-readable media comprising instructions for execu-
2 tion in one or more processors for executing uniprocessor (UP) coded workloads in a
3 multiprocessor (MP) computer system without having to rewrite the UP-coded work-
4 loads' code, ~~the method comprising the steps:~~

5 organizing the UP-coded workloads into one or more concurrency groups,
6 wherein UP-coded workloads in the same concurrency group are not permitted to execute
7 concurrently with one another in the MP computer system;

scheduling first and second execution vehicles that respectively execute on different processors in the MP computer system at substantially the same time;

acquiring a first concurrency group by the first execution vehicle and a second concurrency group by the second execution vehicle; and

executing UP-coded workloads in the first concurrency group through the first execution vehicle at substantially the same time as UP-coded workloads in the second concurrency group are executed through the second execution vehicle.

26. (Original) The computer-readable media according to claim 25, wherein the UP-coded workloads are UP-coded threads, and the first and second execution vehicles are first and second processes.

27. (Original) The computer-readable media according to claim 25, wherein the UP-coded workloads are messages, and the first and second execution vehicles are first and second threads.

28. (Currently Amended) A method for executing workloads in a multiprocessor (MP) computer system, the method comprising the steps:

organizing the workloads into one or more concurrency groups, wherein workloads in the same concurrency group are not permitted to execute concurrently with one another in the MP computer system;

scheduling first and second execution vehicles that respectively execute on different processors in the MP computer system at substantially the same time;

acquiring a first concurrency group by the first execution vehicle and a second concurrency group by the second execution vehicle; and

executing workloads in the first concurrency group through the first execution vehicle at substantially the same time as workloads in the second concurrency group are executed through the second execution vehicle.

1 29. (Original) The method according to claim 28, wherein the step of acquiring the first
2 and second concurrency groups further comprises:

3 dequeuing from a concurrency-group run queue a first concurrency-group data
4 structure associated with the first concurrency group; and

5 dequeuing from the concurrency-group run queue a second concurrency-group
6 data structure associated with the second concurrency group.

1 30. (Original) The method according to claim 29, further comprising:

2 setting a first CG flag in the first concurrency-group data structure to a value indi-
3 cating that the first concurrency group is in a running state; and

4 setting a second CG flag in the second concurrency-group data structure to a
5 value indicating that the second concurrency group is in a running state.

1 31. (Original) The method according to claim 29, further comprising:

2 appending workloads enqueued on a first current queue in the first concurrency-
3 group data structure onto a first active queue in the first concurrency-group data struc-
4 ture; and

5 appending workloads enqueued on a second current queue in the second concu-
6 rency-group data structure onto a second active queue in the second concurrency-group
7 data structure.

1 32. (Original) The method according to claim 31, further comprising:

2 dequeuing workloads in the first and second concurrency groups from the first
3 and second active queues, respectively; and

4 executing the dequeued workloads to completion.

1 33. (Original) The method according to claim 30, further comprising:

2 in response to the first execution vehicle finishing execution of the workloads in
3 the first concurrency group, the first execution vehicle performing the steps:

- 4 A) if at least one workload in the first concurrency group is executable:
5 (i) setting the value of the first CG flag to a value indicat-
6 ing that the first concurrency group is in a queued state;
7 (ii) re-enqueueing the first concurrency-group data struc-
8 ture onto the concurrency-group run queue;
9 B) if there are not any workloads in the first concurrency group that are
10 executable, setting the first CG flag to a value indicating that the first concurrency
11 group is in a suspended state;
12 C) dequeuing from the concurrency-group run queue a third concur-
13 rency-group data structure associated with a third concurrency group; and
14 D) setting a third CG flag in the third concurrency-group data structure to
15 a value indicating that the third concurrency group is in a running state.

1 Please add new claims 34 *et al.*

1 34. (New) A method, comprising:

2 organizing a plurality of workloads into a plurality of concurrency groups,
3 wherein each workload in the same concurrency group are not permitted to execute con-
4 currently with another workload in a microprocessor (MP) computer system;
5 scheduling a plurality of execution vehicles that respectively execute on different
6 processors in the MP computer system at substantially the same time;
7 acquiring by each execution vehicle of the plurality of execution vehicles a con-
8 currency group from the plurality of concurrency groups; and
9 executing workloads in the plurality of concurrency groups through the plurality
10 of execution vehicles at substantially the same time.

1 35. (New) The method according to claim 34, wherein the workloads are uniprocessor
2 (UP) coded threads, and the plurality of vehicles are processes.

1 36. (New) The method according to claim 34, wherein the workloads are messages, and
2 the plurality of vehicles are first and second threads.

1 37. (New) The method according to claim 34, wherein the workloads are uniprocessor
2 (UP) coded workloads.